

Exo 1

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Q1. According to Frama-C RTE 4 runtime errors could happen:

1. a memory error when executing $T[x]=y$ if $x < 0$
2. a memory error when executing $T[x]=y$ if $x \geq 5$
3. an integer overflow when executing $x=x+1$
4. an integer overflow when executing $y=y+100$

Q2. Running Frama-C EVA we get the following results:

- EVA is not able to prove that error 2 won't occur
- EVA is not able to prove that error 4 won't occur

The 2 other errors are discharged (they will definitely not occur)

Q3.

Running a VSA by hand we get the following results at the entry of each basic block:

* without widening, the loop is unrolled up to termination:

(note that variable y is *not* constrained by condition of block B1)

Entering B0

$x=bot, y=bot$

Entering B1

$x=[0,6], y=[0,600]$

Entering B2

$x=[0,5], y=[0,600]$

Entering B3

$x=[1,5], y=[0,600]$

Entering B4

$x=[0,6], y=[0,600]$

Entering B5

$x=[6,6], y=[0,600]$

We get the same conclusion than Frama-C for error 2, but not for error 4 (no integer overflow when incrementing y)

* with widening/narrowing, variables are set to $+infty$ after one iteration:

(note that variable y is *not* narrowed by condition of block B1)

Entering B0

$x=bot, y=bot$

Entering B1

$x=[0,6], y=[0,+infty]$

Entering B2

$x=[0,5], y=[0,+infty]$

Entering B3

$x=[1,5], y=[0,+infty]$

Entering B4

$x=[0,6], y=[0,+infty]$

Entering B5

$x=[6,6], y=[0,+infty]$

We get the same conclusions than Frama-C.

* According to the normal program execution, error 1 may really occur at runtime,

but error 2 will not occur. This last error is therefore a false positive.

(Frama-C is not able to discharge because it cannot catch the implicit relation between x and y ,

which is caught without using widening/narrowing).

Q4. For $N=1000$ we would get:

```
Entering B0
  x=bot, y=bot
Entering B1
  x=[0,1001], y=[0,+infty]
Entering B2
  x=[0,1000], y=[0,+infty]
Entering B3
  x=[1,999], y=[0,+infty]
Entering B4
  x=[0,1000], y=[0,+infty]
Entering B5
  x=[1001,1001], y=[0,+infty]
```

Here error 2 is discharged, because within B3 we have $x < 1000$.
However, error 4 is not discharged, and it is still a false positive ...

Q5.

We want to check under which conditions on N we would get a buffer overflow at line 11.

To do so we need to express a constraint on N (considered as a symbolic value) telling whether line 11 can be executed with $x < 0$ or $x \geq N$.

In practice, since the value of N impacts the number of loop iterations, we need to enumerate several values of N (since each of them leads to a different path predicate for reaching line 11).

For instance, we should consider the following constraints:

- no iteration
 $N=0$ and $x=0$ and $x < N+1$ and $x \% 2 = 1$ and $(x < 0$ or $x \geq N)$,
 which is not satisfiable
- 1 iteration
 $N=1$ and $x=0$ and $x < N+1$ and $x \% 2 = 1$ and $(x < 0$ or $x \geq N)$ and
 $x_1 = x+1$ and $x_1 < N+1$ and $x_1 \% 2 = 1$ and $(x_1 < 0$ or $x_1 \geq N)$,
 which is satisfiable since $x_1 = 1$.
- 2 iterations
 $N=2$ and $x=0$ and $x < N+1$ and $x \% 2 = 1$ and $(x < 0$ or $x \geq N)$ and
 $x_1 = x+1$ and $x_1 < N+1$ and $x_1 \% 2 = 1$ and $(x_1 < 0$ or $x_1 \geq N)$ and
 $x_2 = x_1+1$ and $x_2 < N+1$ and $x_2 \% 2 = 1$ and $(x_2 < 0$ or $x_2 \geq N)$
 which is not satisfiable
- etc.

In conclusion the best we can do here is to check a *finite* set of constraints (corresponding to each numbers of iterations), and we will get an error whenever N is odd.

But we won't be able to conclude for *any* value of N ...