

# Software security, secure programming

## Lecture 1: introduction

Master M2 Cybersecurity

Academic Year 2024 - 2025

# Who are we ?

## Teaching staff

- ▶ Laurent Mounier (UGA)
- ▶ research within Verimag Lab (PACS team)
- ▶ research focus: formal verification, code analysis, compilation techniques, language semantics ... and **(software) security** !

## Attendees

- ▶ Master M2 CySec students

→ various skills, background and interests ...

# Agenda

## Part 1: an overview of software security and secure programming

- ▶ ~ 7 weeks (21 hours)
- ▶ classes on **wednesday** (2pm - 5pm)

## Part 2: some tools and techniques for software security

- ▶ ~ 6 weeks (18 hours)
- ▶ class on **tuesday** (2pm - 5pm)

→ includes lectures, training exercises, labs . . .

# Examination rules

The rules of the game ...

## Assignments

- ▶  $M_1$ : a written assignment (duration=1h, mid-November)
- ▶  $M_2$ : (short) reports on some lab sessions
- ▶  $M_3$ : final written exam (duration=2h, end of January)

## Mark computation (3 ECTS)

$$M = (0.66 \times M_1 + 0.33 \times M_2) + (0.5 \times M_3)$$

## Course user manual

### An (on-going) course web page on **Moodle** ...

`https://im2ag-moodle.univ-grenoble-alpes.fr/course/view.php?id=545`

- ▶ course schedule and materials (slides, past exams, etc.)
- ▶ weekly, reading suggestions, to complete the lecture
- ▶ other background reading/browsing advices ...

### During the classes ...

Alternation between lectures, written exercices, lab exercises ...

...but no “formal” lectures → questions & discussions always welcome !

heterogeneous audience + open topics ⇒ **high interactivity level !**

# Prerequisites

Ideally ...

This course is concerned with:

## Programming languages

- ▶ at least one (classical) imperative language:  
C or C++, Java, maybe Python ...
- ▶ some notions on compilation & (informal) language semantics

## What happens behind the curtain

Some notions about:

- ▶ assembly code (x86, others ? ...)
- ▶ runtime memory layout (stack, heap)

# Outline

Some practical information

What **software security** is (not) about ?

About software security

## The context: computer system security . . .

**Question 1:** what is a “computer system”, or an **execution platform** ?

Many possible incarnations, e.g.:

- ▶ (classical) computer: mainframe, server, desktop, laptop, etc.
- ▶ mobile device: phone, tablets, audio/video player, etc.  
... up to IoT, smart cards, ...
- ▶ embedded (networked) systems: inside a car, a plane, a washing-machine, etc.
- ▶ cloud/remote computing, virtual execution environment
- ▶ but also industrial networks (Scada), ... etc.
- ▶ and certainly many more !

→ 2 main characteristics:

- ▶ include hardware + **software**
- ▶ open/connected to the **outside world** . . .



## The context: computer system security ... (ct'd)

### Question 2: what does mean **security** ?

- ▶ a set of general **security** properties: CIA  
Confidentiality, Integrity, Availability (+ Non Repudiation + Anonymity + ...)
- ▶ concerns the **running** software + the whole **execution platform**  
(other users, shared resources and data, peripherals, network, etc.)
- ▶ depends on an **intruder model**  
→ there is an “external actor”<sup>1</sup> with an **attack objective** in mind, and able to elaborate a dedicated strategy to achieve it (≠ hazards)  
↔ something beyond **safety** and **fault-tolerance**

→ A possible definition:

- ▶ fonctionnal properties = what the system should do
- ▶ security properties = what it should **not allow** w.r.t the intruder model ...

**Rk:** fonctionnal properties do matter for “security-oriented” software (firewalls, etc.)!

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<sup>1</sup>could be the user, or the **execution platform itself!**

## Example 1: password authentication

Is this code “secure” ?

```
boolean verify (char[] input, char[] passwd , byte len) {
    // No more than triesLeft attempts
    if (triesLeft < 0) return false ; // no authentication
    // Main comparison
    for (short i=0; i <= len; i++)
        if (input[i] != passwd[i]) {
            triesLeft-- ;
            return false ; // no authentication
        }
    // Comparison is successful
    triesLeft = maxTries ;
    return true ; // authentication is successful
}
```

functional property:

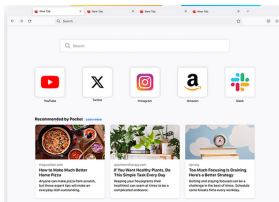
$$\text{verify}(\text{input}, \text{passwd}, \text{len}) \Leftrightarrow \text{input}[0..\text{len}] = \text{passwd}[0..\text{len}]$$

What do we want to protect ? Against what ?

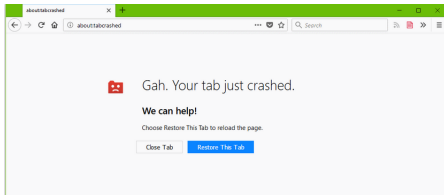
- ▶ confidentiality of `passwd`, information leakage ?
- ▶ no unexpected runtime behaviour
- ▶ code integrity, etc.

## Example 2: web browser

Unavoidable applications, key functionalities, routinely used . . .



But, quite often:



Is it a simple **functionality issue?**

(no damage, users simply need to restart their browser . . .)

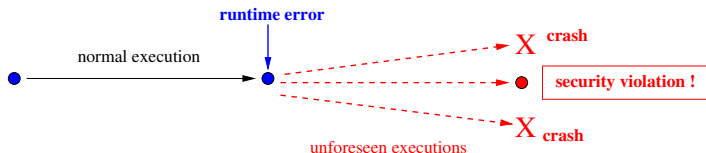
## Why do we need to bother about crashes ?

crash = consequence of an unexpected run-time error

- ▶ not foreseen by the programmer and compiler ...
- ▶ ... and not (always) accurately trapped at runtime

⇒ some part of the execution:

- ▶ may take place **outside the program scope**  
(not following the regular program semantic)
- ▶ but can be **controled/exploited** by an attacker (~ “weird machine”)



⇔ may **break** all security properties ...  
from simple denial-of-service to **arbitrary code execution**

**Rk:** may also happen **silently** (without any crash !)

## Some (not standardized) definitions ...

**Bug:** an error (or defect/ flaw/ failure) introduced in a SW, either

- ▶ at the specification / design / algorithmic level
- ▶ at the programming / coding level
- ▶ or even by the compiler (or any other pgm transformation tools) ...

**Vulnerability:** a weakness (for instance a bug !) that opens a “security breach”

- ▶ **non exploitable** vulnerabilities: there is no (known !) way for an attacker to use this bug to corrupt the system
- ▶ **exploitable** vulnerabilities: this bug can be used to elaborate an attack (i.e., write an **exploit**)
- ▶ **0-day** vulnerabilities: yet unpublished (hence **not patched** !)

**Exploit:** a concrete attacker behavior allowing to:

1. trigger a (sequence of) vulnerability(-ies)
2. leading to a security property violation

Ex: a single program input, or a complex sequence of interactions with the target program and/or its execution environment ...

# Software vulnerability examples

## Case 1 (not so common ...)

### Functional property not provided by a security-oriented component

- ▶ lack of encryption, too weak crypto-system,
- ▶ no (strong enough) authentication mechanism,
- ▶ bad firewall configuration, too weak access control enforcement rules,
- ▶ etc.

## Case 2 (the vast majority !)

### Insecure coding practice in (any!) software component/application

- ▶ improper input validation  $\rightsquigarrow$  SQL or code injection, XSS, etc.
- ▶ insecure shared resource management (file system, network)
- ▶ information leakage (lack of data encapsulation, side channels)
- ▶ exploitable coding errors (memory access, arithmetic overflows, etc.)
- ▶ etc.

⇒ **Sleeping bombs**



# The intruder model

## Who/what is the attacker ?

- ▶ a malicious external user, interacting via regular input sources e.g., keyboard, network (man-in-the-middle), etc.
- ▶ a malicious external “observer”, interacting via side channels (execution time, power consumption)
- ▶ another application running on the same platform interacting through shared resources like caches, processor elements, etc.
- ▶ the execution platform itself (e.g., when compromised !)

## What is he/she/it able to do ?

At low level:

- ▶ unexpected memory read (data or code)
- ▶ unexpected memory write (data or code)

⇒ **powerful enough for**

- ▶ information disclosure
- ▶ unexpected/arbitrary code execution
- ▶ privilege elevation, etc.

## Example: smartphone attack surface



Credits [BT2019]



# Outline

Some practical information

What software security is (not) about ?

About software security

## Some evidences regarding cyber (un)-security

So many examples of successful computer system attacks:

- ▶ the “**famous ones**”: (at least one per year !)  
Morris worm, Stuxnet, Heartbleed, WannaCry, Spectre, Log4j, etc.
- ▶ the never-ending records of “**cyber-attacks**” against large organizations (private companies, public structures)
- ▶ a public database of **CVEs** (Common Vulnerabilities and Exposures)  
Numbers of CVEs per year
- ▶ etc.

Why ? Who can we blame for that ??

- ▶ ∄ well defined recipe to build secure cyber systems in the large
- ▶ permanent trade-off between **efficiency** and **safety/security**:
  - ▶ HW and micro-architectures (**sharing** is everywhere !)
  - ▶ operating systems
  - ▶ programming languages and applications
  - ▶ coding and software engineering techniques

## But, what about **software** security ?

Software is **greatly involved** in “computer system security”:

- ▶ it plays a major role in **enforcing security properties**:  
crypto, authentication protocols, intrusion detection, firewall, etc.
- ▶ but it is also a major source of **security problems**<sup>2</sup> . . .  
“90 percent of security incidents result from exploits against defects in software” ( U.S. DHS)

→ SW is clearly one of the **weakest links** in the security chain!

### Why ???

- ▶ we do not do very well how to write **secure** SW  
we do not even know how to write **correct** SW!
- ▶ behavioral properties can't be validated on a (large) SW  
impossible by hand, untractable with a machine
- ▶ programming languages not designed for security enforcement  
most of them contain numerous traps and pitfalls
- ▶ programmers feel not (so much) concerned with security  
security not get enough attention in programming/SE courses
- ▶ heterogenous and nomad applications favor unsecure SW  
remote execution, mobile code, plugins, reflection, etc.

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<sup>2</sup>outside security related code!

# Some concrete CVE examples: back to the browsers . . .

## 🚩 CVE-2022-26485 Detail

### Description

Removing an XSLT parameter during processing could have lead to an exploitable use-after-free. We have had reports of attacks in the wild abusing this flaw. This vulnerability affects Firefox < 97.0.2, Firefox ESR < 91.6.1, Firefox for Android < 97.3.0, Thunderbird < 91.6.2, and Focus < 97.3.0.

### Metrics

CVSS Version 4.0 CVSS Version 3.x CVSS Version 2.0

*NVD enrichment efforts reference publicly available information to associate vector strings. CVSS information contributed by other sources is also displayed.*

#### CVSS 3.x Severity and Vector Strings:



NIST: NVD

Base Score: **8.8 HIGH**

Vector: CVSS:3.1/AV:N/AC:L/PR:N/UI:R/S:U/C:H/I:H/A:H

## 🚩 CVE-2024-29944 Detail

### AWAITING ANALYSIS

This vulnerability is currently awaiting analysis.

### Description

An attacker was able to inject an event handler into a privileged object that would allow arbitrary JavaScript execution in the parent process. Note: This vulnerability affects Desktop Firefox only, it does not affect mobile versions of Firefox. This vulnerability affects Firefox < 124.0.1 and Firefox ESR < 115.9.1.

### Metrics

CVSS Version 4.0 CVSS Version 3.x CVSS Version 2.0

*NVD enrichment efforts reference publicly available information to associate vector strings. CVSS information contributed by other sources is also displayed.*

#### CVSS 4.0 Severity and Vector Strings:



NIST: NVD

**N/A**

NVD assessment not yet provided.

See the [online discussions](#) . . .

# A highly critical recent CVE example (Trojan Horse)

## 🚩 CVE-2024-3094 Detail

### MODIFIED

This vulnerability has been modified since it was last analyzed by the NVD. It is awaiting reanalysis which may result in further changes to the information provided.

## Description

Malicious code was discovered in the upstream tarballs of xz, starting with version 5.6.0. Through a series of complex obfuscations, the liblzma build process extracts a prebuilt object file from a disguised test file existing in the source code, which is then used to modify specific functions in the liblzma code. This results in a modified liblzma library that can be used by any software linked against this library, intercepting and modifying the data interaction with this library.

## Metrics

CVSS Version 4.0

CVSS Version 3.x

CVSS Version 2.0

*NVD enrichment efforts reference publicly available information to associate vector strings. CVSS information contributed by other sources is also displayed.*

### CVSS 3.x Severity and Vector Strings:



**CNA:** Red Hat, Inc.

**Base Score:** 10.0 CRITICAL

**Vector:** CVSS:3.1/AV:N/AC:L/PR:N/UI:N/S:C/C:H/I:H/A:H

(see the [Pentest-Tools blog](#))

And [more CVEs](#) are still coming !

## Some evidences regarding software (un)-security (ct'd)

An increasing activity in the “defender side” as well ...

- ▶ all the daily security patches (for OS, basic applications, etc.)
- ▶ companies and experts specialized in software security  
code audit, search for 0days, malware detection & analysis, etc.  
“bug bounties” [<https://zerodium.com/program.html>]
- ▶ some important research efforts  
from the main software editors (e.g., MicroSoft, Google, etc)  
from the academia (conferences) and independent “ethical hackers” (blogs, etc.)
- ▶ software verification tools editors start addressing security issues  
e.g.: dedicated static analyser features
- ▶ international cooperation for vulnerability disclosure and classification  
e.g.: CERT, CVE/CWE catalogue, vulnerability databases
- ▶ government agencies to promote & control SW security  
e.g.: ANSSI, ENISA, Darpa “Grand Challenge”, etc.
- ▶ national/european/international regulations, norms and standards  
e.g.: RGPD, NIS-2, Cyber Resilience Act, ISO 27001, IEC 62443

## Couter-measures and protections (examples)

Several existing mechanisms to **enforce** SW security

- ▶ at the programming level:
  - ▶ disclosed vulnerabilities → language weaknesses databases  
↳ secure coding patterns and libraries
  - ▶ aggressive compiler options + code instrumentation  
↳ early detection of unsecure code
  
- ▶ at the OS level:
  - ▶ sandboxing
  - ▶ address space randomization
  - ▶ non executable memory zones
  - ▶ etc.
  
- ▶ at the hardware level:
  - ▶ Trusted Platform Modules (TPM)
  - ▶ secure crypto-processor
  - ▶ CPU tracking mechanisms (e.g., Intel Processor Trace)
  - ▶ etc.

## Techniques and tools for assessing SW security

Several existing mechanisms to **evaluate** SW security

- ▶ **code review** ...
- ▶ **fuzzing**:
  - ▶ run the code with “unexpected” inputs → **pgm crashes**
  - ▶ (tedious) manual check to find **exploitable** vulns ...
- ▶ **(smart) testing**:  
coverage-oriented pgm exploration techniques  
(genetic algorithms, dynamic-symbolic executions, etc.)  
+ code instrumentation to detect (low-level) vulnerabilities
- ▶ **static analysis**: approximate the code behavior to detect **potential** vulns  
(~ code optimization techniques)

### In practice:

- ▶ only **the binary code** is **always available** and **useful** ...
- ▶ **combinations** of all these techniques ...
- ▶ **exploitability analysis** still challenging ...



## Course objectives (for the part 1)

Understand the **root causes** of common weaknesses in SW security

- ▶ at the programming language level
- ▶ at the execution platform level

→ helps to better choose (or deal with) a programming language

Learn some methods and techniques to **build more secure SW**:

- ▶ programming techniques:  
languages, coding patterns, etc.
- ▶ validation techniques:  
what can(not) bring existing tools ?
- ▶ counter-measures and protection mechanisms

## Course agenda

See

<https://im2ag-moodle.univ-grenoble-alpes.fr/course/view.php?id=545>

## Credits:

- ▶ E. Poll (Radboud University)
- ▶ M. Payer (Purdue University)
- ▶ E. Jaeger, O. Levillain and P. Chifflier (ANSSI)