Three-Address Code

- Or "**TAC**"
- The IR that you will be using for the final programming project.
- High-level assembly where each operation has at most three operands.
- Uses explicit runtime stack for function calls.
- Uses vtables for dynamic dispatch.

```
int x;
int y;

int x2 = x * x;
int y2 = y * y;
int r2 = x2 + y2;
```

```
int x;
int y;

int x2 = x * x;
int y2 = y * y;
int r2 = x2 + y2;
```

```
x2 = x * x;

y2 = y * y;

r2 = x2 + y2;
```

```
int a;
int b;
int c;
int d;

a = b + c + d;
b = a * a + b * b;
```

```
int a;
int b;
int c;
int d;

a = b + c + d;
b = a * a + b * b;
```

```
_t0 = b + c;
a = _t0 + d;
_t1 = a * a;
_t2 = b * b;
b = _t1 + _t2;
```

```
int a;
int b;
int c;
int d;

a = b + c + d;
b = a * a + b * b;
```

```
_t0 = b + c;
a = _t0 + d;
_t1 = a * a;
_t2 = b * b;
b = _t1 + _t2;
```

Temporary Variables

- The "three" in "three-address code" refers to the number of operands in any instruction.
- Evaluating an expression with more than three subexpressions requires the introduction of temporary variables.
- This is actually a lot easier than you might think; we'll see how to do it later on.

```
int a;
int b;
a = 5 + 2 * b;
```

```
int a;
int b;
a = 5 + 2 * b;
```

```
int a;
int b;
a = 5 + 2 * b;
```

TAC allows for instructions with two operands.

Simple TAC Instructions

Variable assignment allows assignments of the form

```
var = constant;
var<sub>1</sub> = var<sub>2</sub>;
var<sub>1</sub> = var<sub>2</sub> op var<sub>3</sub>;
var<sub>1</sub> = constant op var<sub>2</sub>;
var<sub>1</sub> = var<sub>2</sub> op constant;
var = constant<sub>1</sub> op constant<sub>2</sub>;
```

- Permitted operators are +, -, *, /, %.
- How would you compile y = -x;?

Simple TAC Instructions

 Variable assignment allows assignments of the form

```
var = constant;
• var_1 = var_2;
var<sub>1</sub> = var<sub>2</sub> op var<sub>3</sub>;
• var<sub>1</sub> = constant op var<sub>2</sub>;
var<sub>1</sub> = var<sub>2</sub> op constant;
var = constant, op constant,;
```

- Permitted operators are +, -, *, /, %.
- How would you compile y = -x;?

$$y = 0 - x; y = -1 * x;$$

One More with bools

```
int x;
int y;
bool b1;
bool b2;
bool b3;

b1 = x + x < y
b2 = x + x == y
b3 = x + x > y
```

One More with bools

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int x;
int y;
bool b1;
bool b2;
bool b3;

b1 = x + x < y
b2 = x + x == y
b3 = x + x > y
```

```
t0 = x + x;
t1 = y;
\overline{b}1 = t0 < t1;
 t2 = x + x;
t3 = y;
\overline{b}2 = t2 == t3;
 t4 = x + x;
t5 = y;
b3 = t5 < t4;
```

TAC with bools

- Boolean variables are represented as integers that have zero or nonzero values.
- In addition to the arithmetic operator,
 TAC supports <, ==, ||, and &&.
- How might you compile $b = (x \le y)$?

TAC with bools

- Boolean variables are represented as integers that have zero or nonzero values.
- In addition to the arithmetic operator,
 TAC supports <, ==, ||, and &&.
- How might you compile $b = (x \le y)$?

```
int x;
int y;
int z;
if (x < y)
   z = x;
else
   z = y;
z = z * z;
```

```
int x;
int y;
int z;

if (x < y)
   z = x;
else
   z = y;</pre>
```

```
__t0 = x < y;
IfZ __t0 Goto __L0;
z = x;
Goto __L1;
__L0:
z = y;
__L1:
z = z * z;
```

```
int x;
int y;
int z;

if (x < y)
    z = x;
else
    z = y;</pre>
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```
__t0 = x < y;
IfZ __t0 Goto __L0;
z = x;
Goto __L1;
__L0:
z = y;
__L1:
z = z * z;
```

Labels

- TAC allows for **named labels** indicating particular points in the code that can be jumped to.
- There are two control flow instructions:
 - Goto label;
 - IfZ value Goto label;
- Note that Ifz is always paired with Goto.

```
int x;
int y;

while (x < y) {
    x = x * 2;
}

y = x;</pre>
```

```
int x;
int y;

while (x < y) {
    x = x * 2;
}

y = x;</pre>
```

```
_L0:
_t0 = x < y;
IfZ _t0 Goto _L1;
x = x * 2;
Goto _L0;
_L1:
_y = x;
```